Collaboration in social networks: Incentives and topology

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Luca Dall'Asta, Matteo Marsili, and Paolo Pin Collaboration in social networks

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The puzzle of cooperation

Why do we see so much cooperation around?

Failed states, why do societies collapse?

Will Euro collapse if Greece drops out?







Much has been written on the emergence of cooperation on networks

- Repeated games, reputation and trust (Myerson 1991)
- Endogenous network games (Vega-Redondo 2007, Jackson 2008, Goyal 2009)
- Repeated games on evolving networks (Ellison 1994, Haag Lagunoff 2006, Vega-Redondo 2006).
- Cooperation in evolutionary games without mutation (Boyd 1999, Hofbauer Sigmund 2003, Poncela et al 2010)
- Repeated games and punishment on specific structures (Eshel et al 1998, Haag Lagunoff 2007, Fainmesser 2009, Karlan et al 2009)
- Focus here: social network = pattern of repeated interactions repeated interaction = forward looking behavior collaboration = incentives + credibility of threats How difficult is this in large games on complex structures?

Outline

- The prisoners dilemma
- Collaboration in repeated interaction: 2 players
 - Collaboration is supported by credible threats of punishment
- Collaboration in N players games on a network: Local contribution game
 - Conditional collaboration has to be reciprocal and limited to a subset of neighbors
- How does collaboration depend on incentives and topology?
 - Collaborative equilibria are subgraphs of the social network
- The complexity of collaboration:
 - Counting collaborative equilibria with message passing
- Conclusions

Defection is the only possible outcome in one shot prisoner's dilemma



N players on graph G=(N,L) Each player either cooperates (C) or defects (D) with all neighbors

Payoff: 1 for each neighbor that collaborates minus X_i (=cost of collaboration)

All D ($s_i=0$) is the only Nash equilibrium



 $u_i(s_i, s_{-i}) = -X_i s_i + \sum s_j$

 $s_i = 0, 1$

 $j \in \partial_i$

N=2: When the game is played many times cooperation is possible, among other things

Strategies become plans of actions, decided at time 0, to optimize future payoffs

$U_i = (1 - d) \sum_{t=0}^{\infty} d^t u_i \left(s_i^{(t)}, s_{-i}^{(t)} \right), \quad d \in [0, 1]$

- Cooperation under trigger strategies T: T= {start with C;
 - C as long as opponent plays C,
 - D forever, if opponent plays D}

If d is large enough, (T, T) is a Nash equilibrium

- Folk's theorem: many other outcomes can be supported as a Nash equilibrium
- I d=1 in what follows



But threats should be credible

N=3

- Is it credible that 1 and 2 punish 3?
- Not if $u_1(C,C,D) > u_1(D,D,D)$!
- Players need to condition C only to a subset of their neighbors
- If i conditions on j, j should condition on i
- Emergent heterogeneity





On trees, Nash equilibria are subtrees

- Given an undirected tree G=(N,L) $k_i = |\partial_i| = degree of node i$ $m_i = smallest integer larger than X_i$ $c_i = number of collaborators in <math>\partial_i$
- Any collection of disjoint undirected subgraphs $\Gamma = (V, \Lambda)$ of G is a collaborative equilibrium where all $i \in V$ cooperate conditionally to neighbors in Γ and $|\partial_i \cap \Lambda| = m_i$
- Incentives: $i \in V$ $c_i X_i \ge c_i m_i \Rightarrow m_i \ge X_i$
- Reciprocity: i,j∈V, if j does not punish i
 ⇒ i should not punish j when j defects
- Credibility: $i,k \in V, (i,k) \in \Lambda$ if k defects $c_i - 1 - X_i < c_i - m_i \Rightarrow m_i < X_i + 1$

$u_i(s_i, s_{-i}) = c_i - X_i s_i$

On generic graphs cascades of defection make things more complex

- Indirect defections: As a result of the defection of j∈∂_i other neighbors k∈∂_i may also defect because of loops
- A collection of disjoint undirected subgraphs Γ=(V,Λ) of G is a collaborative equilibrium where all i∈V cooperate conditionally to neighbors in Γ and |∂_i ∩ Λ|=m_i provided
 i) the indirect effects caused by the defection of all j∈∂_i ∩ Λ have the same consequence of the defection of i itself.
- i) holds provided removing i from V does not disconnect Γ
 - Works on trees, for dimers and loops, for the complete graph
 - Likely works on random graphs and on dense graphs

counter example

Nash equilibria on random graphs

Do NE exist? How many? How hard is it to find them?

Circuits: Marinari, Monasson, Semerjian 2006 q-regular subgraphs: Pretti, Weigt 2006

Counting NE by message passing

• $x_{i \rightarrow j} = 1$ if i conditions C on j, $x_{i \rightarrow j} = 0$ otherwise

- there are m_i-1 k $\in \partial_i/j$ with $x_{k \rightarrow i}=1 \Rightarrow x_{i \rightarrow j}=1$
- $m_i k \in \partial_i / j$ with $x_{k \rightarrow i} = 1 \Rightarrow x_{i \rightarrow j} = 0$
- no $k \in \partial_i/j$ with $x_{k \rightarrow i} = 1 \implies x_{i \rightarrow j} = 0$

Marginals:

$\mu_{i \to j} = P\{i \in V, i \text{ punishes } j\}$

Circuits: Marinari, Monasson, Semerjian 2006 q-regular subgraphs: Pretti, Weigt 2006

Message passing equations:

$$\begin{split} \mu_{i \to j} &= \frac{e^{-\epsilon} Z_{N_i \setminus j \to i}^{m_i - 1}}{Z_{N_i \setminus j \to i}^0 + e^{-\epsilon} Z_{N_i \setminus j \to i}^{m_i - 1} + e^{-\epsilon} Z_{N_i \setminus j \to i}^{m_i}} \\ Z_{V \to i}^q &= \sum_{U \subseteq V} \mathbb{I}_{|U| = q} \prod_{j \in U} \mu_{j \to i} \prod_{k \in V/U} (1 - \mu_{k \to i}) \\ P\{i \in C\} &= \frac{e^{-\epsilon} Z_{N_i \to i}^{\gamma_i}}{Z_{N_i \to i}^0 + e^{-\epsilon} Z_{N_i \to i}^{\gamma_i}} \\ P\{i \in \Gamma_j\} &= \frac{\mu_{i \to j} \mu_{j \to i}}{\mu_{i \to j} \mu_{j \to i} + (1 - \mu_{i \to j})(1 - \mu_{j \to i})}. \end{split}$$

• Fixed point \Rightarrow number of subgraphs (entropy)

Circuits: Marinari, Monasson, Semerjian 2006 q-regular subgraphs: Pretti, Weigt 2006

k

Regular random graphs: dimers (mi=1)

Exponentially many NE's

• ρ = fraction of cooperators s(ρ) = log(number of NE| ρ)/N

NE $\exists \forall \rho \in [0,1]$

Regular random graphs: circuits (mi=2)

Exponentially many NE's

• ρ = fraction of cooperators s(ρ) = log(number of NE| ρ)/N

NE $\exists \forall \rho \in [0,1]$

(Marinari, Monasson, Semerjian 2006)

Regular random graphs: m_i=3

- Exponentially many NE's
- NE $\nexists \forall \rho < \rho_c$
- NE are non-local and fragile

Heterogeneous random graphs

Scale free: hubs collaborate more likely than spokes

Erdös-Rényi: E[k]=4

$X_i = x k_i$

Assortative networks are more conducive to collaboration

Scale free network $P(k) \sim k^{-2.5}$ $X_i = xk_i, x=0.1$

Conclusions: Theory

Collaboration in repeated prisoners dilemma as a graph theoretical problem:

- 1- make sure enough neighbors collaborate
- 2- not credible to monitor more neighbors
- 3- checks should be reciprocal
- If incentives to defect (x)
 - is small then cooperation is easy
 - is large

i) collaboration requires critical mass

- ii) Nash equilibria are fragile
- iii) effect of defection are non-local
- Topology: Collaboration is easier on
 i) trees

ii) densely connected graphs

Collaboration is harder on networks which can be disconnected (e.g. quasi 1d graphs)

Conclusion: Empirical evidence

- Individuals condition collaboration on that of others (Fishbacher, Gachter, Fehr 2001)
- Weak and strong ties (Granovetter, 1973) Individuals do not condition collaboration to all contacts, not even to all those who collaborate, only to a subset of them
- Critical mass theory of collective action (Oliver, Marwell 1993) If the cost of collaboration is large enough, collaborative equilibria only arise if a finite fraction of agents participate
- Collaboration easier in dense networks (Kirchkamp, Nagel 2007; Cassar 2007)
- More connected agents are more likely to collaborate (Cassar 2007)
- Collaboration is not contagious (Suri, Watts 2011)
 The more of my contacts are engaged in conditional collaboration with others, the less likely I am to find neighbors with whom to collaborate conditionally